

Formal Logic

Lecture 9: The Semantics of Predicate Logic (Part II)

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Validity and other Notions in L_2

Logical truth in L_2

- A sentence ϕ in L_2 is *logically true* IFF it is true under any interpretation.

NB: Bergmann et al. (2014) call it ‘quantificationally true’.

Examples:

$$(\forall x)(Px \rightarrow Px)$$

$$(\exists x)(Px \vee \neg Px)$$

- We cannot show that a sentence is logically true via one or even some interpretations. We need all of them.
- We can show that a sentence is *not* logically true by coming up with one interpretation that makes the sentence false.

Contradiction in L_2

- A sentence ϕ in L_2 is *logically false* IFF it is false under any interpretation.

NB: Bergmann et al. (2014) call it ‘quantificationally false’.

Examples:

$$\neg(\exists x)(Px \vee \neg Px)$$

$$(\forall x) Px \ \& \ \neg Pa$$

- We cannot show that a sentence is logically false via one or even some interpretations. We need all of them.
- We can show that a sentence is *not* logically false by coming up with one interpretation that makes the sentence true.

Contingency in L_2

- A sentence ϕ in L_2 is *contingent* if and only if it is neither logically true nor logically false.

NB: Bergmann et al. (2014) call it ‘quant. indeterminate’.

Examples:

$(\exists x)(Px \ \& \ Qx)$

$(\forall y) \text{ Ray}$

- We can show that a sentence is contingent via two interpretations (one where it comes out true and one false).

Validity in L_2

- An argument in L_2 is *valid* IFF there is no interpretation under which the premises are true and the conclusion is false.

NB: Bergmann et al. (2014) call it ‘quantificationally valid’.

Examples:

$(\forall x) Px$	entails	Pa
$\{(\forall x)(Px \rightarrow Qa), (\exists x)Px\}$	entails	Qa

- We cannot show that an argument is valid via one or even some interpretations. We need all of them.
- We can show that an argument is invalid via 1 interpretation that makes the premises true and the conclusion false.

Consistency in L_2

- A set of sentences in L_2 is *consistent* IFF there is at least one interpretation on which all the members of the set are true.

NB: Bergmann et al. (2014) call them ‘quant. consistent’.

Examples:

$\{(\exists x)Hx, \neg Hc\}$

$\{(\forall y)Ray, \neg Rba \vee (\exists z) \neg Raz\}$

- We cannot show that a set of sentences is inconsistent via one or even some interpretations. We need all of them.
- We can show that a set of sentences is consistent via one interpretation that makes all sentences true.

Logical equivalence in L_2

- Two sentences ϕ, ψ in L_2 are *logically equivalent* IFF there is no interpretation on which ϕ, ψ have different truth-values.

NB: Bergmann et al. (2014) call them ‘quant. equivalent’.

Examples:

$$\begin{array}{lcl} (\forall x) \neg Qx & \text{and} & \neg(\exists x) Qx \\ (\exists x)(Ax \ \& \ \neg Bx) & \text{and} & \neg(\forall x)(Ax \rightarrow Bx) \end{array}$$

- We cannot show that a pair of sentences is equivalent via one or even some interpretations. We need all of them.
- We can show that a pair is *not* logically equivalent via 1 interpretation where one sentence is true and the other false.

The undecidability of L_2

- Alonzo Church (1936) proved a fundamental limit to first-order predicate logic.
 - There is no decision procedure that determines for every group of L_2 sentences whether they are:
 - * logically true, logically false or contingent
 - * consistent or inconsistent
 - * valid or invalid
 - * equivalent or inequivalent
- NB:** In L_1 , there is such a procedure, namely truth-tables.
- **Decision procedure:** A mechanical method that shows for all cases & in a finite number of steps whether a property holds.

Exercise Set #5

Exercise 4.1

- **Instructions:** Determine whether the following expressions are formulae of L_2 and say which of those are also sentences of L_2 .

Exercise 4.1

Formula, sentence or neither?

- (i) $(\forall x)(P_1x \rightarrow Qy)$
- (ii) $(\exists x) \neg(\neg\neg(\exists y)Py \ \& \ \neg\neg\neg\neg\neg\neg Rxa)$
- (iii) P
- (iv) $(\forall x)(\exists y)(\exists z)(R_{25}xyz)$
- (v) $(\forall x)(\exists x) Qxx$
- (vi) $\neg(\neg((\exists x)Px \ \& \ (\exists y)Qy))$
- (vii) $(\forall x)(\exists y)(Pxy \ \& \ Px) \vee Qxyx$

Exercise 4.1: Solution

Formula, sentence or neither?

(i) $(\forall x) (P_1x \rightarrow Qy)$

(ii) $(\exists x) \neg(\neg\neg(\exists y)Py \ \& \ \neg\neg\neg\neg\neg Rxa)$

(iii) P

(iv) $(\forall x)(\exists y)(\exists z) (R_{25}xyz)$

(v) $(\forall x)(\exists x) Qxx$

(vi) $\neg(\neg((\exists x)Px \ \& \ (\exists y)Qy))$

(vii) $(\forall x)((\exists y)(Pxy \ \& \ Px) \vee Qxyx)$

Formula (y is free)

Sentence

Sentence

Neither (brackets on R)

Sentence

Neither (outer brackets)

Formula (2nd y is free)

Exercise 4.3

- **Instructions:** Provide L_2 -formalisations for the following English sentences.

Exercise 4.3

English Sentences:

- (i) London is big and ugly.
- (ii) Culham is a large village.
- (iii) A city has a city hall.
- (iv) Material objects are divisible.
- (v) Tom owns at least one car.
- (vi) Tom owns at least one car and
he won't sell it.
- (vii) One man has visited every country.

Exercise 4.3: Solution

English Sentences:

- (i) London is big and ugly.
- (ii) Culham is a large village.
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- (vii) One man has visited every country.

Formalised:

$Ba \& Ua$

Lc

$(\forall x)(Cx \rightarrow (\exists y)(Pxy \& Hy))$

$(\forall x)(Mx \rightarrow Dx)$

$(\exists x)(Otx \& Cx)$

$(\exists x)((Otx \& Cx) \& \neg Stx)$

$(\exists x)(Mx \& (\forall y)(Cy \rightarrow Vxy))$

NB: Other readings sometimes possible, e.g. one means one in vii.

Exercise 4.4

- **Instructions:** Translate the L_2 -sentences below into English using the following dictionary.

Exercise 4.4

- Translate the L_2 -sentences below into English using the following dictionary.

a: Tom

P^1 : ... is a person

Q^1 : ... acts freely

L_2 Sentences:

(i) Qa

(ii) $(Qa \vee \neg Pa)$

(iii) $(\forall x) (Px \rightarrow Qx)$

(iv) $(\forall x) (Px \leftrightarrow Qx)$

(v) $\neg \exists z_1 Qz_1$

Exercise 4.4: Solution

- Translate the L_2 -sentences below into English using the following dictionary.

a: Tom

P^1 : ... is a person

Q^1 : ... acts freely

L_2 Sentences:

(i) Qa

(ii) $(Qa \vee \neg Pa)$

(iii) $(\forall x) (Px \rightarrow Qx)$

(iv) $(\forall x) (Px \leftrightarrow Qx)$

(v) $\neg \exists z_1 Qz_1$

English Sentences:

Tom acts freely.

Tom acts freely or Tom is not a person.

Every person acts freely.

Every person acts freely and vice-versa.

Nothing acts freely.

Exercise 5.1

- **Instructions:** Consider an L_2 -structure S ... Are the following sentences true or false in this structure? Sketch proofs...

Exercise 5.1

- Consider an L_2 -structure S ... Are the following sentences true or false in this structure? Sketch proofs...

UD: $\{1, 2, 3\}$, $P: \{2\}$, $R: \{<1, 2>, <2, 3>, <1, 3>\}$, $a: 1$, $b: 3$

(i) Pa

(ii) Rab

(iii) Rba

(iv) $Rab \leftrightarrow Rba$

(v) $Rbb \vee (\neg Pa \ \& \ \neg Raa)$

(vi) $(\exists x)Rax$

(vii) $(\exists x)(Rax \ \& \ Rxb)$

(viii) $Pb \vee (\exists x)Rxx$

(ix) $(\forall x)(\exists y)Rxy$

Exercise 5.1: Solution

- Consider an L_2 -structure S ... Are the following sentences true or false in this structure? Sketch proofs...

UD: $\{1, 2, 3\}$, $P: \{2\}$, $R: \langle 1, 2 \rangle, \langle 2, 3 \rangle, \langle 1, 3 \rangle$, $a: 1$, $b: 3$

- | | |
|--|---|
| (i) Pa | False as $1 \notin P$ |
| (ii) Rab | True as $\langle 1, 3 \rangle \in R$ |
| (iii) Rba | False as $\langle 3, 1 \rangle \notin R$ |
| (iv) $Rab \leftrightarrow Rba$ | False as Rab is true and Rba false |
| (v) $Rbb \vee (\neg Pa \ \& \ \neg Raa)$ | True as $1 \notin P$ and $\langle 1, 1 \rangle \notin R$ |
| (vi) $(\exists x)Rax$ | True, e.g. $\langle 1, 2 \rangle \in R$ |
| (vii) $(\exists x)(Rax \ \& \ Rxb)$ | True as $\langle 1, 2 \rangle \in R$ and $\langle 2, 3 \rangle \in R$ |
| (viii) $Pb \vee (\exists x)Rxx$ | False as neither $3 \in P$ nor $\langle x, x \rangle \in R$ |
| (ix) $(\forall x)(\exists y)Rxy$ | False as $\langle 3, y \rangle \notin R$ |

The End